

CSE 420

Computer Architecture I

Brief Review

Computer Organization & Assembly Language

Prof. Michel A. Kinsy

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Software Mechanics for Bridging

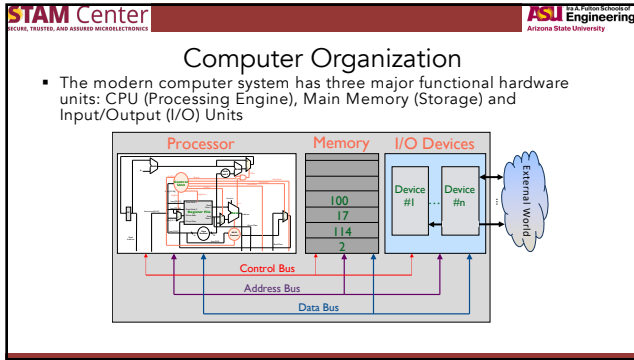
- The Art of Abstraction

Application
Algorithm
Programming Language
Operating System/Virtual Machine
Instruction Set Architecture (ISA)
Microarchitecture
Register-Transfer Level (RTL)
Circuits
Devices
Physics

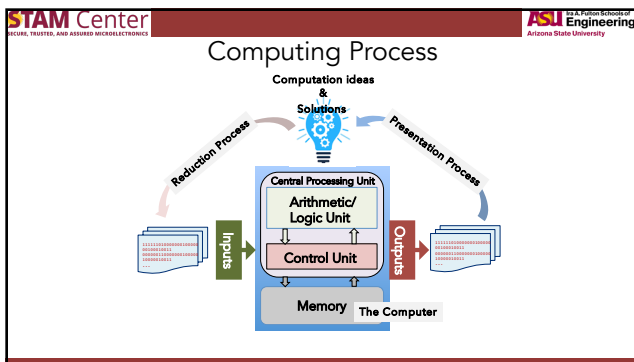
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Another View of the Abstraction

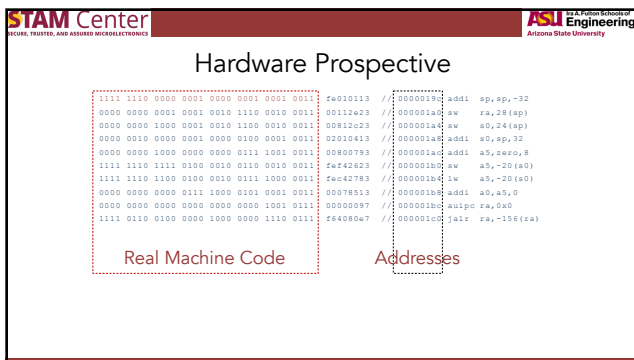
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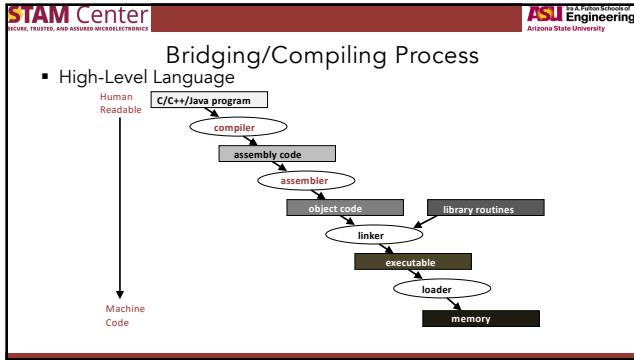
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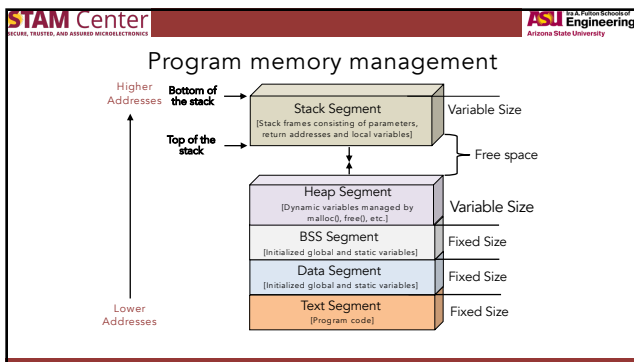
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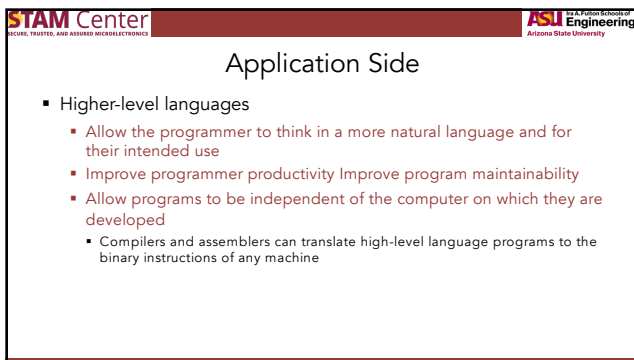
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Application Side

- Higher-level languages
 - Allow the programmer to think in a more natural language and for their intended use
 - Improve programmer productivity Improve program maintainability
 - Allow programs to be independent of the computer on which they are developed
 - Emergence of optimizing compilers that produce very efficient assembly code
 - As a result, very little programming is done today at the assembler level

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System Software Side

- System software
 - Operating system – supervising program that interfaces the user's program with the hardware (e.g., Linux, MacOS, Windows)
 - Handles basic input and output operations
 - Allocates storage and memory
 - Provides for protected sharing among multiple applications
 - Compiler – translate programs written in a high-level language (e.g., C, Java) into instructions that the hardware can execute

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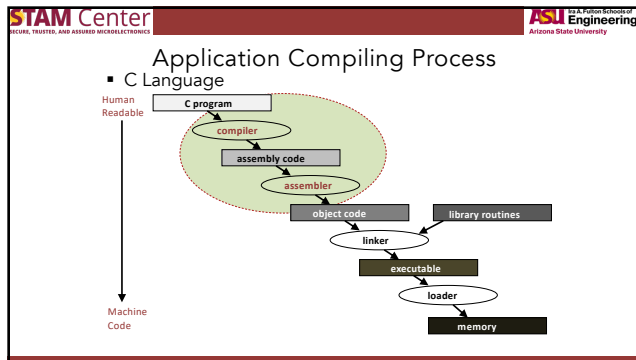
Application Compiling Process

- C Language

```

graph TD
    A[C program] --> B((compiler))
    B --> C[assembly code]
    C --> D((assembler))
    D --> E[object code]
    E --> F((linker))
    G[library routines] --> F
    F --> H[executable]
    H --> I((loader))
    I --> J[memory]
    J --> K[Machine Code]
    A --> L[Human Readable]
    K --> M[Machine Code]
  
```

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Why is assembly level view?

- To become familiar with the process of compiling a program/application (e.g., C) onto a computer system
- To know what assemblers are and what compilers do
- To understand the computer hardware view of the program/application

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Why is assembly level view?

- To become familiar with the process of compiling a program/application (e.g., C) onto a computer system
- To, then, fully realize why computers are built the way they are
 - In turn, you will gain new insights into how to write better and more efficient code
 - And explore new opportunities in the field of embedded system programming

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Greatest Common Divisor Example

```

int gcd(int a, int b) {
    int tmp;
    if(a < b) {
        tmp = a;
        a = b;
        b = tmp;
    }
    //Find the gcd
    while(b != 0) {
        while(a >= b) {
            a = a - b;
        }
        tmp = a;
        a = b;
        b = tmp;
    }
    return a;
}

```

```

main:
    sd ra,24(sp)
    ...
    call printf
    addi a4,a0,28
    ...
    call scanf
    lw a5,-24(a0)
    lw a4,-28(a0)
    mv a3,a4
    mv a0,a5
    call gcd(int,int)
    mv a5,a0
    sw a5,20(s0)
    ...
    call printf
    addi sp,sp,32
    jr ra

```

From C to assembly, the translation is straightforward

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Hardware Prospective

<pre> 1111 1110 0000 0001 0000 0001 0001 0011 0000 0000 0001 0001 0010 1110 0010 0011 0000 0000 1000 0001 0010 1100 0010 0011 0000 0010 0000 0001 0000 0100 0001 0011 0000 0000 1000 0000 0000 0111 1001 0011 1111 1110 1111 0100 0010 0110 0010 0011 1111 1110 1100 0100 0010 0111 1000 0011 0000 0000 0000 0111 1000 0101 0001 0011 0000 0000 0000 0000 0000 0000 1001 0111 1111 0110 0100 0000 1000 0000 1110 0111 </pre>	<pre> fe010113 // 00000110: addi sp,sp,-32 00112423 // 00000110: sw ra,28(sp) 00812c23 // 00000110: sw a0,24(sp) 02010413 // 00000110: addi a0,sp,32 00800793 // 00000110: addi a5,zero,8 fe42623 // 00000110: sw a5,-20(s0) fec42783 // 00000110: lw a5,-20(s0) 00078513 // 00000110: addi a0,a5,0 00000097 // 00000110: auipc ra,0x0 fe4080e7 // 00000110: jalr ra,-156(ra) </pre>
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Real Machine Code Addresses

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Assembly Code

- Three types of statements in assembly language
 - Typically, one statement per a line
 - 1. Executable assembly instructions
 - Operations to be performed by the processor
 - 2. Pseudo-Instructions and Macros
 - Translated by the assembler into real assembly instructions
 - Simplify the programmer task
 - 3. Assembler Directives
 - Provide information to the assembler while translating a program
 - Used to define segments, allocate memory variables, etc.

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Computer Organization Overview

- The modern digital computer has three major functional hardware units: CPU, Main Memory and Input/Output (I/O) Units

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Assembly Code

- There are 3 main types of assembly instructions
 - Arithmetic**
 - add, sub, mul, sll, srl, and, or, etc.
 - Load/store**
 - lw,sw,lb,sb
 - Conditional – branches**
 - beq, bne, j, jra

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Assembly Code


- There are 3 main types of assembly instructions
 - Arithmetic**
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 - Load/store**
 - lw,sw,lb,sb
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 - beq, bne, j, jra

```

.L5:
    lw a4, -36(a0)
    lw a5, -40(a0)
    beq a4, a5, .L4
    lw a4, -36(a0)
    lw a5, -40(a0)
    sub a5, a4, a5
    sw a5, -36(a0)
    j .L5
.L4:
    lw a5, -36(a0)
      
```

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Assembly Code


- There are 3 main types of assembly instructions
 - Arithmetic
 - add, sub, mul, sll, srl, and, or, etc.
 - Load/store
 - lw,sw,lb, sb
 - Conditional – branches
 - beq, bne, j, jra
- Assembly language instructions have the format:
 - [label:] mnemonic [operands] [#comment]

```

.L2
beqz x1, done    # if(x1 == 0) goto done
            
```

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Assembly Code


- There are 3 main types of assembly instructions
 - Arithmetic
 - add, sub, mul, sll, srl, and, or, etc.
 - Load/store
 - lw,sw,lb, sb
 - Conditional – branches
 - beq, bne, j, jra
- Assembly language instructions have the format:
 - [label:] mnemonic [operands] [#comment]

```

main:
addi sp,sp,-32
sd ra,24(sp)
            
```

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Assembly Code

- There are 3 main types of assembly instructions
- Assembly language instructions have the format:
 - [label:] mnemonic [operands] [#comment]
- Label: (optional)
 - Marks the address of a memory location
 - Typically appear in data and text segments

```

int array [] = {2, 4, 5, 0, 1, 7};
int main(void) {
    int x,y,z;
    x = array[0];
    y = array[1];
    z = array[2];
    ...
}
            
```

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Assembly Code

- There are 3 main types of assembly instructions
- Assembly language instructions have the format:
 - [label:] mnemonic [operands] [#comment]
- Label: (optional)
 - Marks the address of a memory location
 - Typically appear in data and text segments

```

int array [] = {2, 4, 5, 0, 1, 7}; array:
int main(void) {
    int x,y,z;
    x = array[0];
    y = array[1];
    z = array[2];
    ...

    main:
    addi sp,sp,-48
    sw ra,44(sp)
    sw s0,40(sp)
    addi s0,sp,48
    lui a5,%hi(array)
    lw x5,%lo(array)(a5)
    lw x6,4(a5)
    lw x7,8(a5)
    
```

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Assembly Code

- .DATA directive
- .TEXT directive
- .GLOBL directive
 - Declares a symbol as global

```

int array [] = {2, 4, 5, 0, 1, 7};
char name [9];
int main(void) {
    int x,y,z;
    x = array[0];
    y = array[1];
    z = array[2];
    ...

    .globl main
    .type main, @function
    main:
    addi sp,sp,-48
    sw ra,44(sp)
    ...
    
```

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Assembly Code

- .DATA directive
- .TEXT directive
- .GLOBL directive
- .BSS directive
 - The BSS contains variables that are initialized to zero or are explicitly initialized in code

```

int array [] = {2, 4, 5, 0, 1, 7};
char name [9];
int main(void) {
    int x,y,z;
    x = array[0];
    y = array[1];
    z = array[2];
    ...

    .globl name
    .bss
    .size name, 9
    name:
    .zero 9
    .text
    .align 1
    
```

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Assembly Code

- .DATA directive
 - Defines the data segment of a program containing data
 - The program's variables should be defined under this directive
- .TEXT directive
 - Defines the code segment of a program containing instructions
- .GLOBL directive
 - Declares a symbol as global
- .BSS directive
 - The BSS contains variables that are initialized to zero or are explicitly initialized in code

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Assembly Code

```
.LC0:
.string "Enter positive
integer a and b: "
.align 2
.LC1:
.string "%d %d"
.align 2
.LC2:
.string "%d = %d"
.text
.align 1
.globl main
.type main, @function
main:
addi sp, sp, -48
sw ra, 44(sp)
...
```

Directive	Arguments	Description
.byte		6-bit comma separated words (unaligned)
.bword		32-bit comma separated words (unaligned)
.half		16-bit comma separated words (naturally aligned)
.word		32-bit comma separated words (naturally aligned)
.ascz	"string"	emit string (alias for .string)
.string	"string"	emit string
.macro	name arg1 [, argn]	begin macro definition [argname to substitute]
.type	symbol, @function	accepted for source compatibility
...

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Assembly Languages

- Assemblers:
 - Convert mnemonic operation codes to their machine language equivalents
 - Convert symbolic operands to their equivalent machine addresses
 - Build the machine instructions in the proper format
 - Convert the data constants to internal machine representations
 - Write the object program and the assembly listing

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System Calls

- Programs do input/output through system calls
- To obtain services from the operating system
- Using the syscall system services
- Issue the syscall instruction


```

addi a0,a5,%lo(.LC0)
call printf
...
call scanf
lw a5,-36(a0)
...
            
```
- Retrieve return values, if any, from result registers

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Application Compiling Process

- High-level language program (in C)


```

void swap (int array[], int i) {
    int temp;
    temp  = array[i];
    array[i] = array[i+1];
    array[i+1] = temp;
}
            
```
- Assembly language program (for RISC-V)


```

swap:
    addi sp,sp,-48
    ...
    mv a5,a1
    ...
    ld s0,40(sp)
    addi sp,sp,48
    jr ra
            
```

one-to-many

C compiler

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Application Compiling Process

- A compiler is a software program that translates a human-oriented high-level programming language code into computer-oriented machine language

```

graph LR
    A[Source Program  
(C, C++, etc.)] --> B[Compiler]
    B --> C[Target Program  
(RISC-V, MIPS, x86, etc.)]
    B --> D[Error messages]
    E[Input] --> C
    C --> F[Output]
    
```

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Application Compiling Process

- Assembly language program (for RISC-V)


```

swap:
    addi sp,sp,-48
    ...
    mv a5,a1
    ...
    ld s0,40(sp)
    addi sp,sp,48
    jr ra
      
```
- Machine (object, binary) code (for RISC-V)


```

111111010000 00010 000 00010 0010011
000000110000 00010 000 01000 0010011
      
```

one-to-one

assembler

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Application Compiling Process

- Detailed compilation process

Language focused transformations

High-level language

Scanner (lexical analysis)

Parser (syntax analysis)

Semantic Analysis (semantic analysis)

Architecture focused transformations

Code Optimizer

Code Generator

Target language

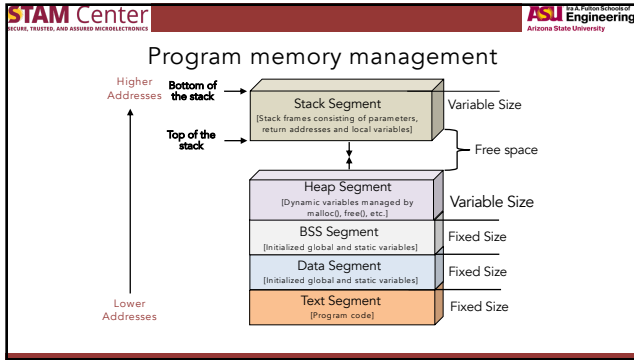
Symbols & Attributes Table
- More on this later when you take a course on compilers

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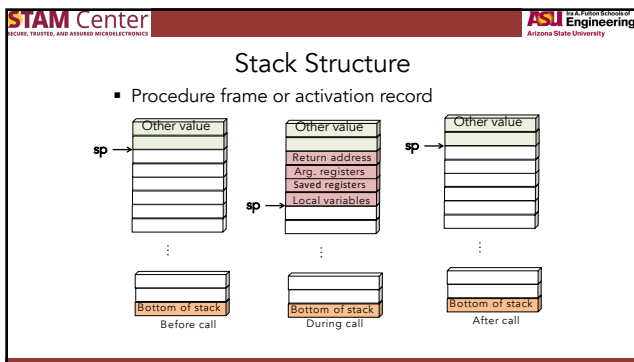
Application Compiling Process

- Symbol Table
 - Identifiers are names of variables, constants, functions, data types, etc.
 - Store information associated with identifiers
 - Information associated with different types of identifiers can be different
 - Information associated with variables are name, type, address, size (for array), etc.

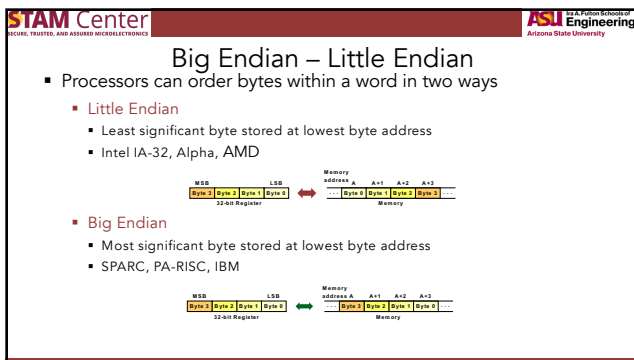
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Big Endian – Little Endian

```

int main(void) {
    int var;           // Integer values
    char *ptr;         // Pointer

    // Assign 'var' and output it in byte order and as a value
    var = 0x12345678;
    ptr = (char *) &var;

    printf("ptr[0] = %02X \n", ptr[0]); // Prints 78
    printf("ptr[1] = %02X \n", ptr[1]); // Prints 56
    printf("ptr[2] = %02X \n", ptr[2]); // Prints 34
    printf("ptr[3] = %02X \n", ptr[3]); // Prints 12

    printf("var = %08X \n", var);       // Prints 12345678
}
    
```

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Big Endian – Little Endian

```

int main(void) {
    int var;           // Integer values
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    // Assign 'var' and output it in byte order and as a value
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    printf("ptr[3] = %02X \n", ptr[3]); // Prints 12

    printf("var = %08X \n", var);       // Prints 12345678
}
    
```

Big Endian	Little Endian
Solaris on SPARC	Windows on Intel
ptr[0] = 12	ptr[0] = 78
ptr[1] = 34	ptr[1] = 56
ptr[2] = 56	ptr[2] = 34
ptr[3] = 78	ptr[3] = 12
var = 12345678	var = 12345678

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Concluding Note

- If you feel the need to learn or refresh some of these foundational concepts, you might consider taking CSE 420 first.

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