

CSE 598

Secure Microkernel Design

Introduction to the fundamentals of computer
architecture using the RISC-V ISA

Prof. Michel A. Kinsky

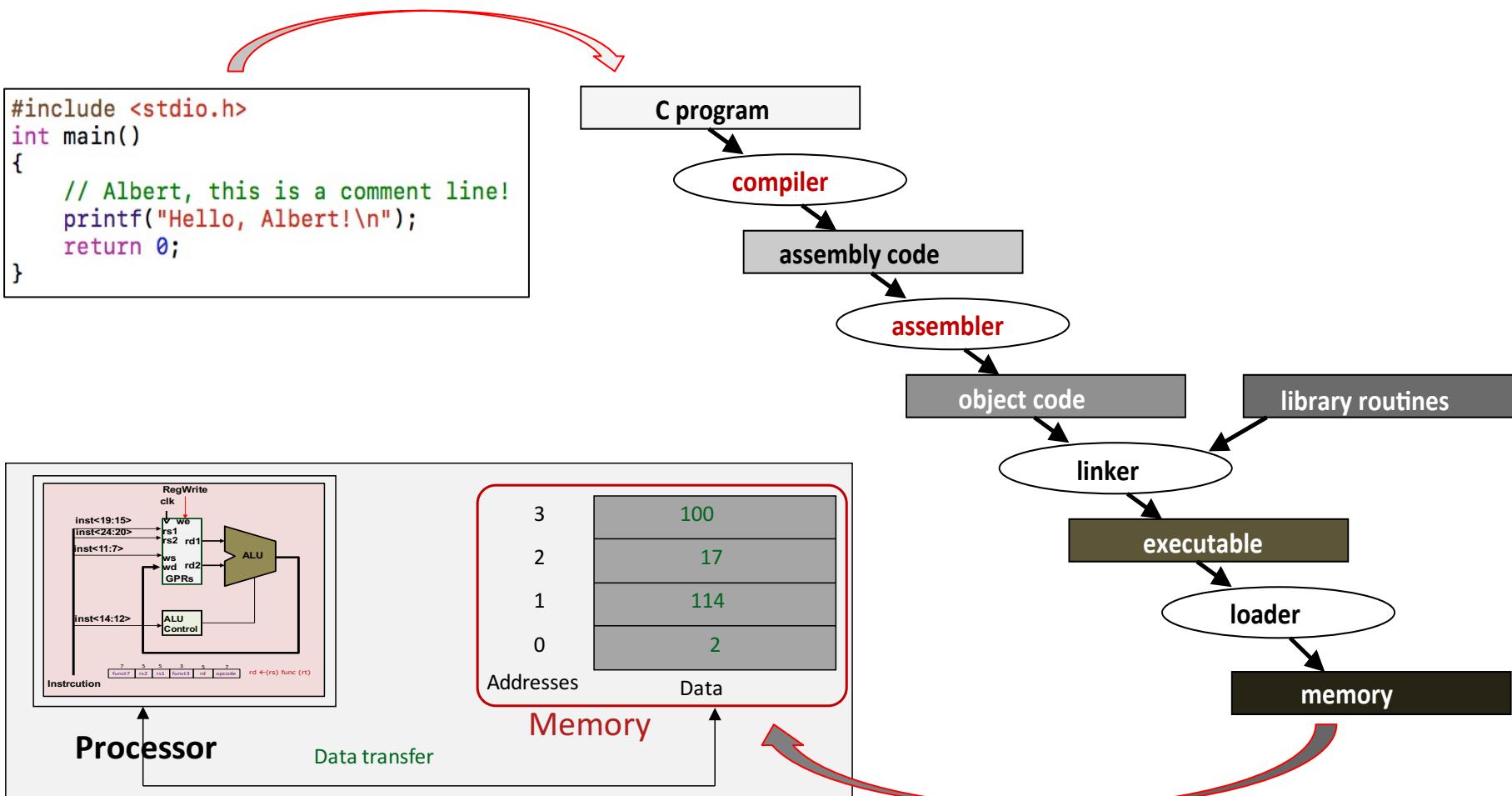
CSE 598

Secure Microkernel Design

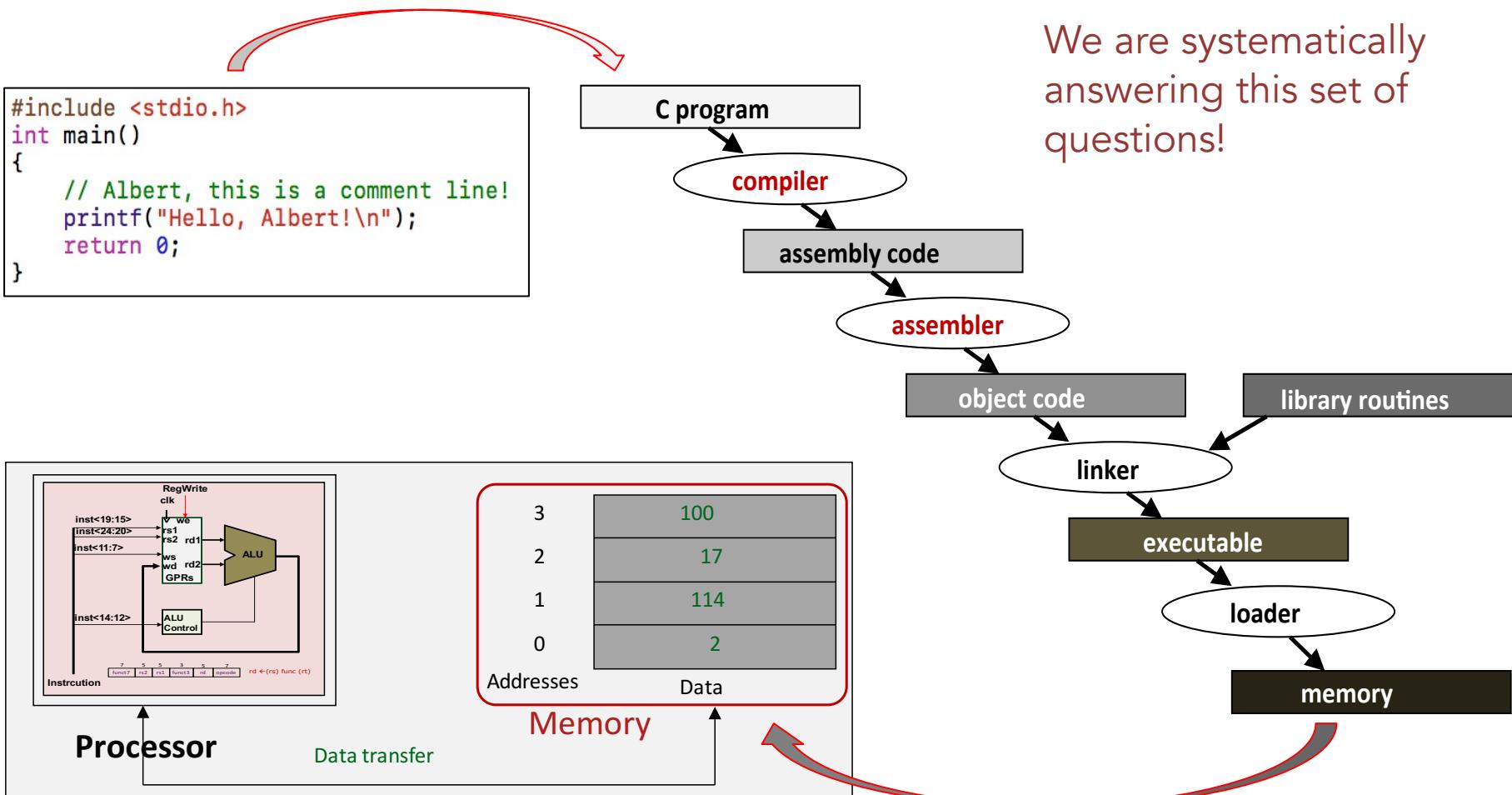
Assembly Language Programming

Prof. Michel A. Kinsky

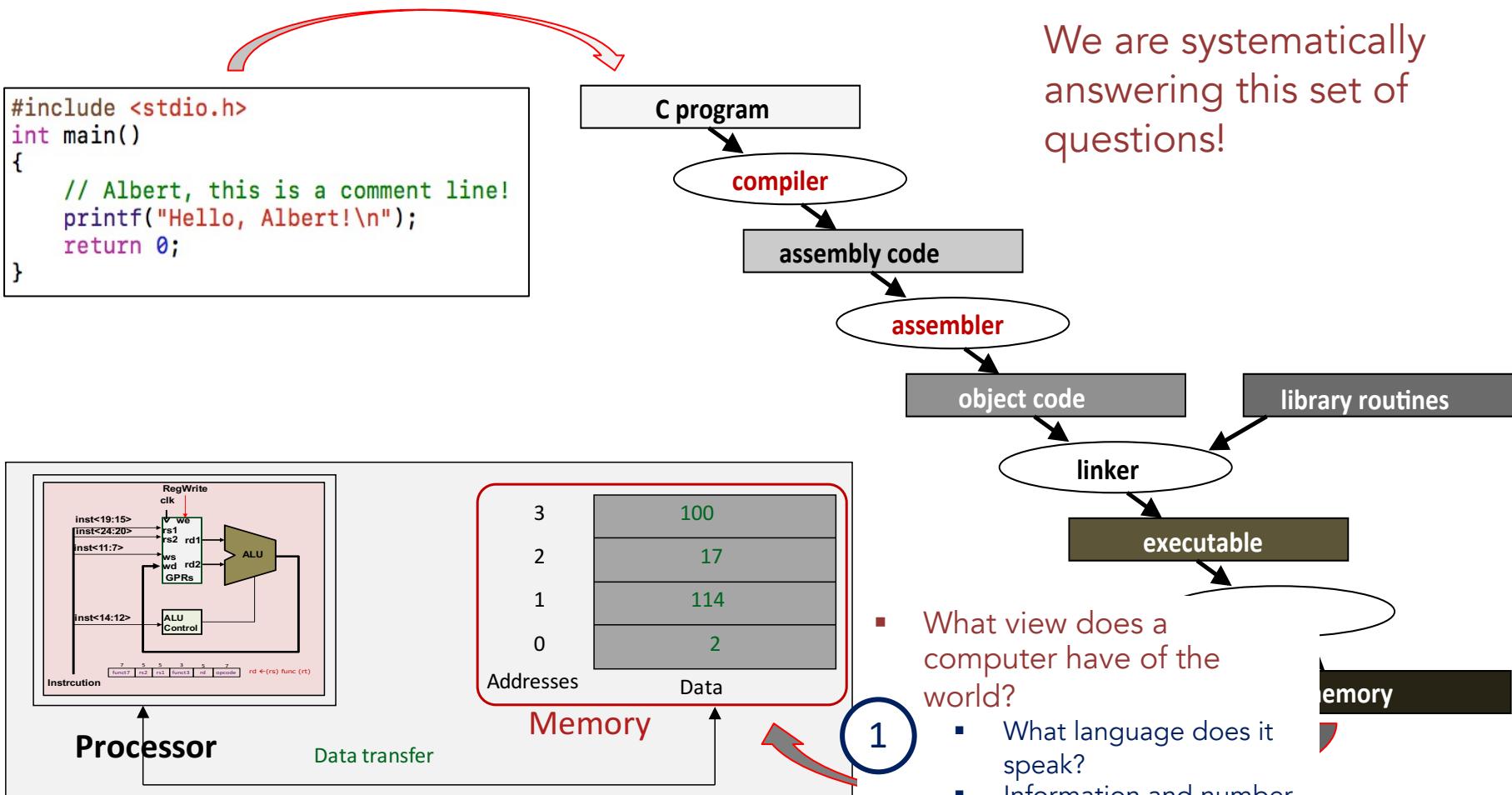
The Big Picture



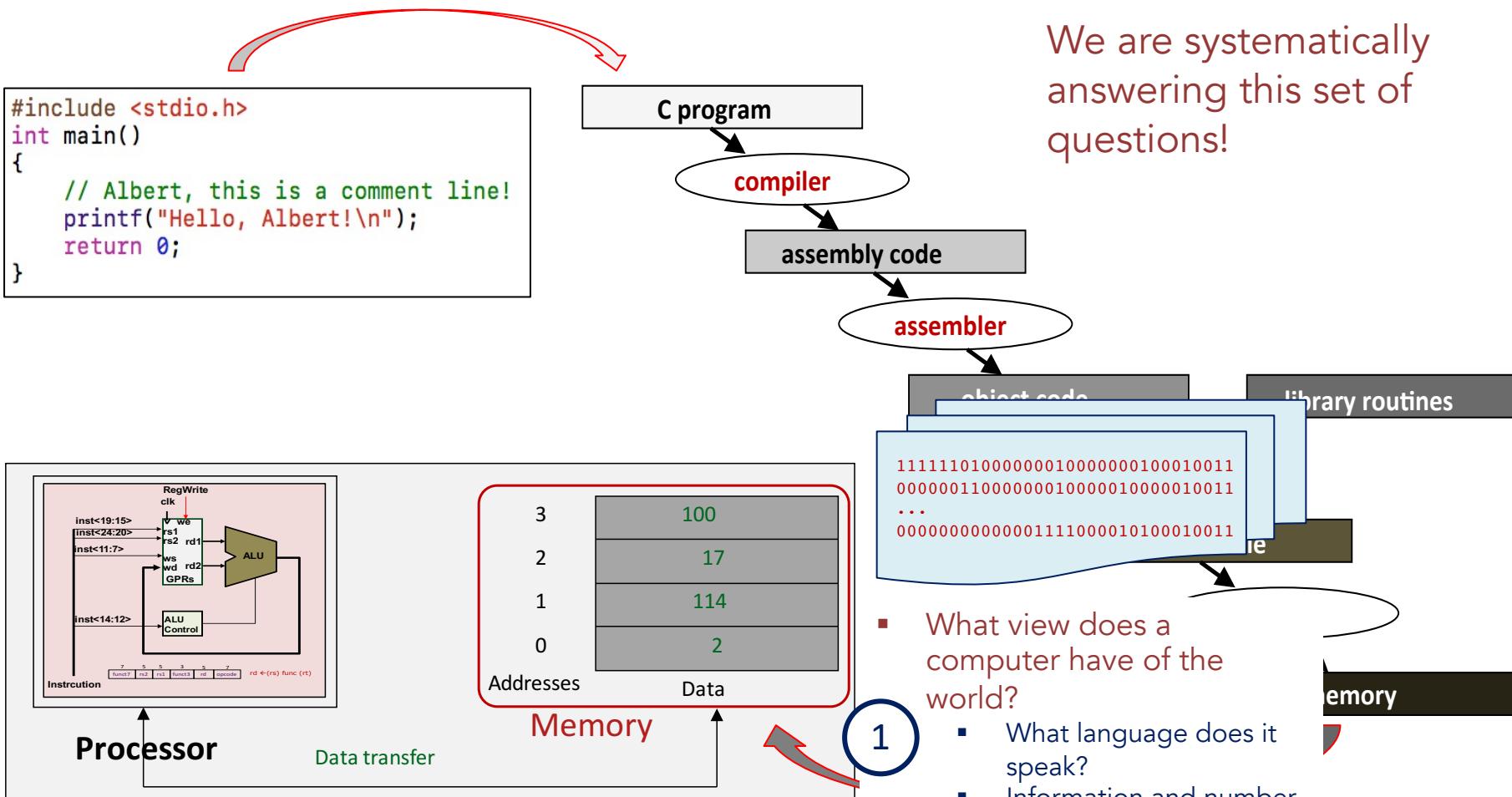
The Big Picture



The Big Picture

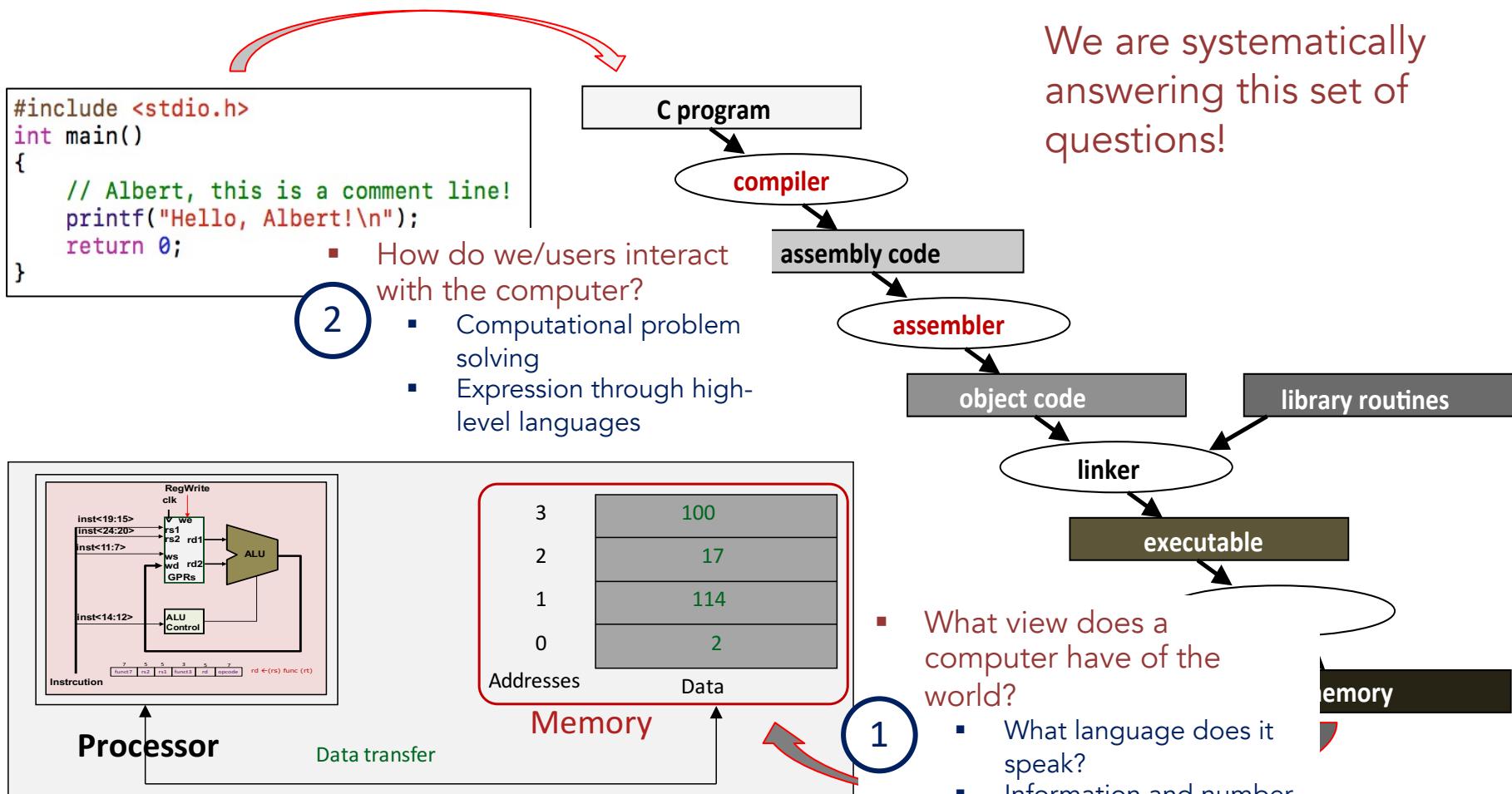


The Big Picture

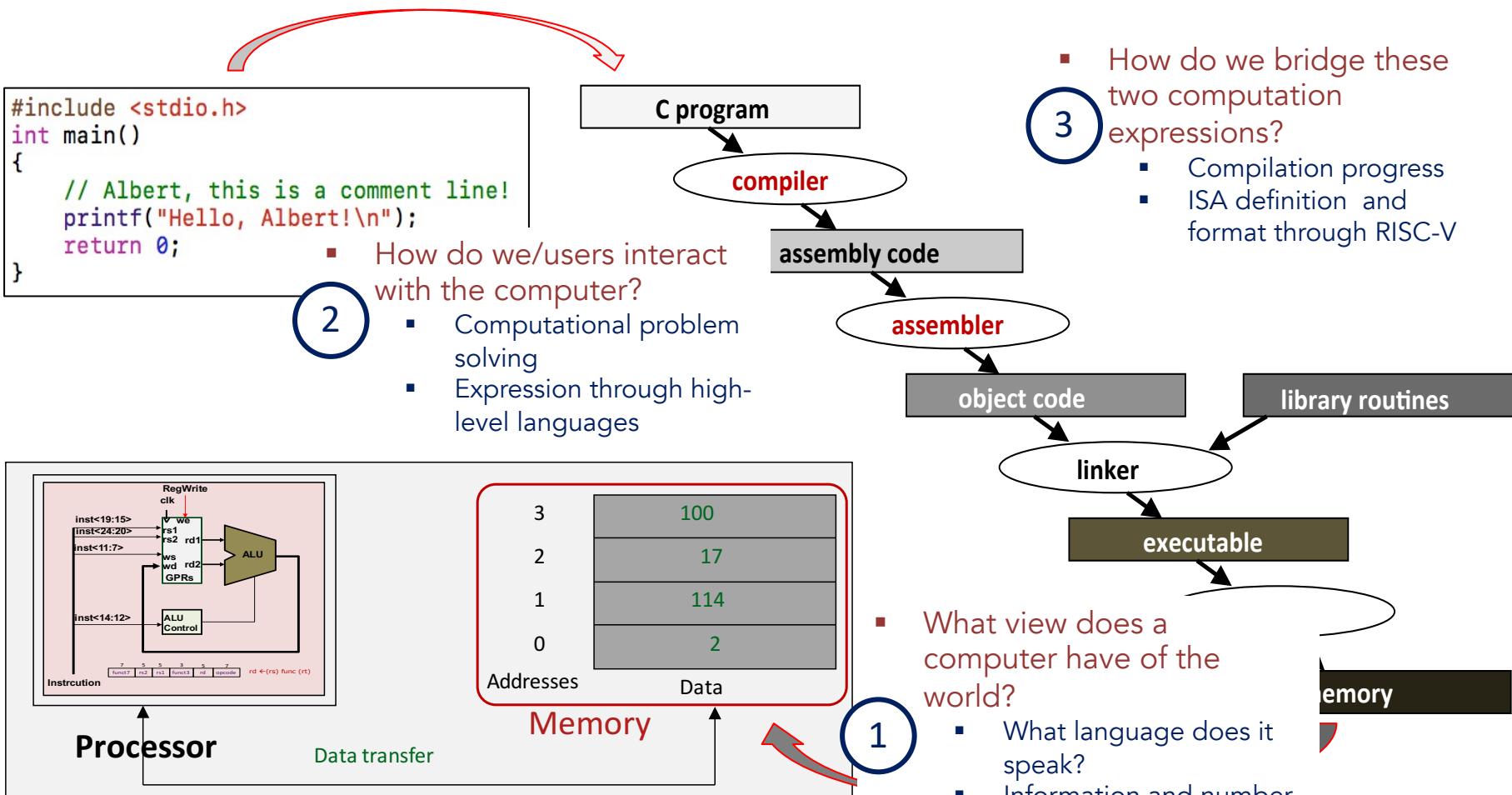


We are systematically answering this set of questions!

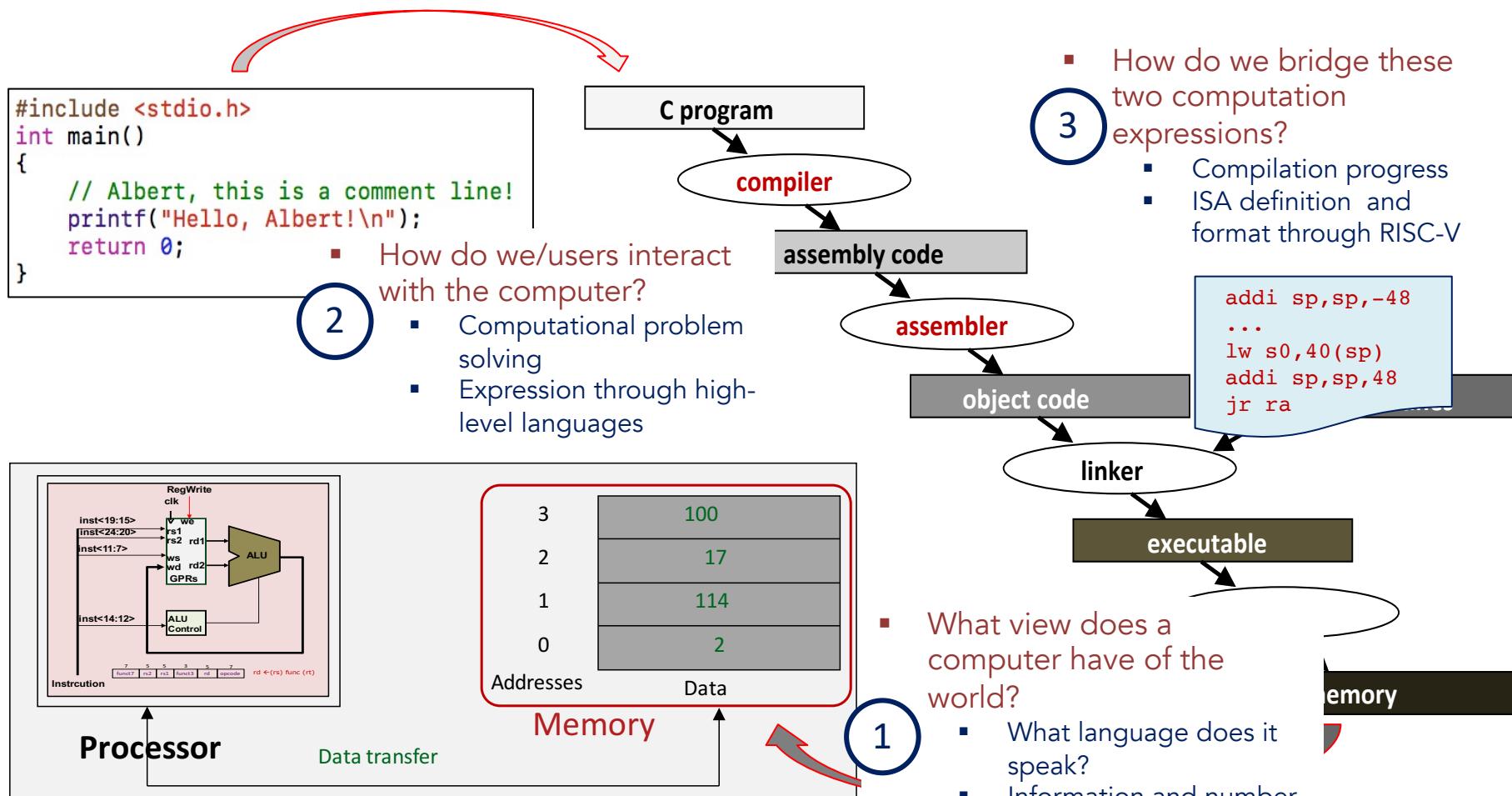
The Big Picture



The Big Picture



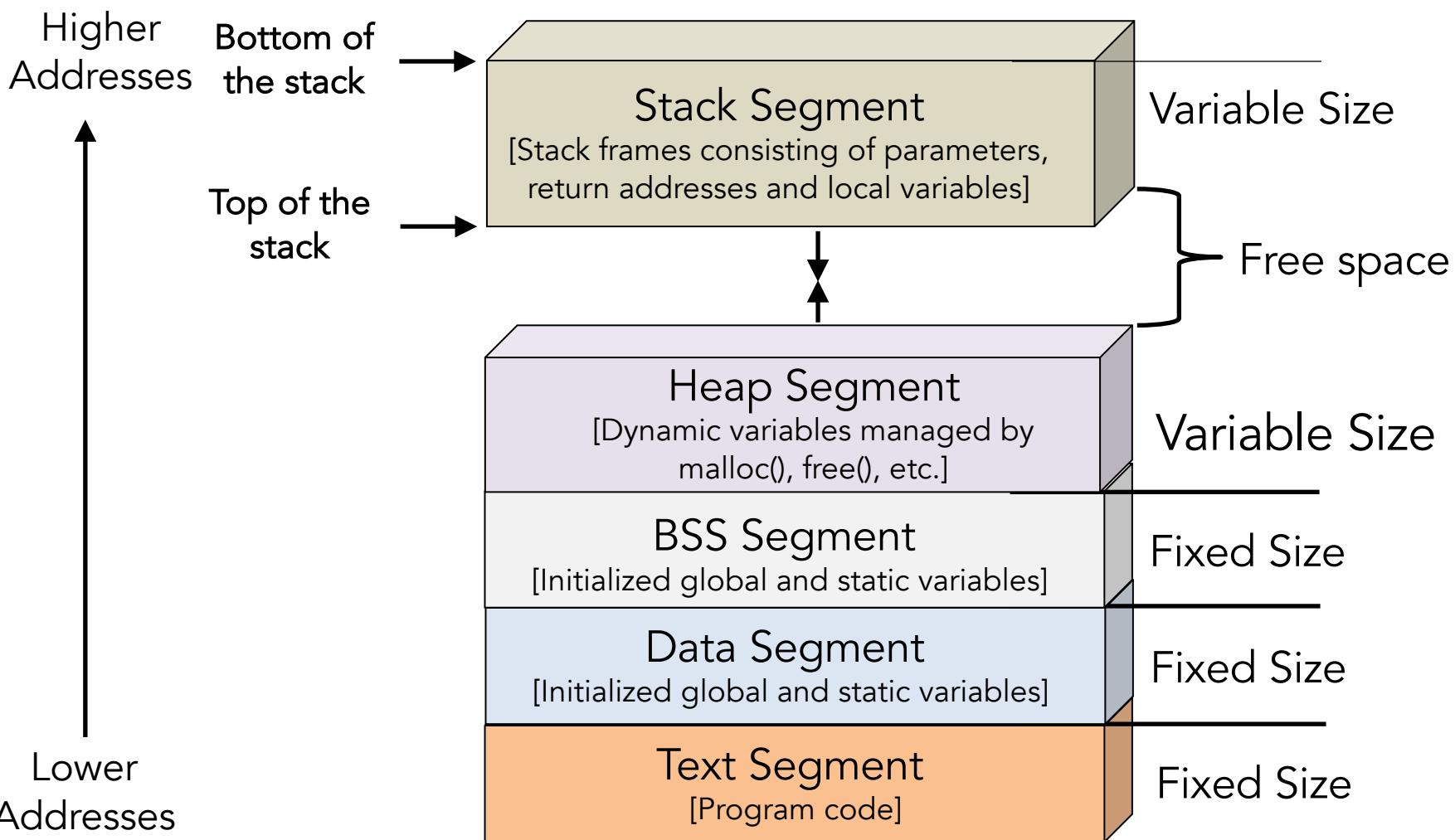
The Big Picture



RISC-V ISA Registers

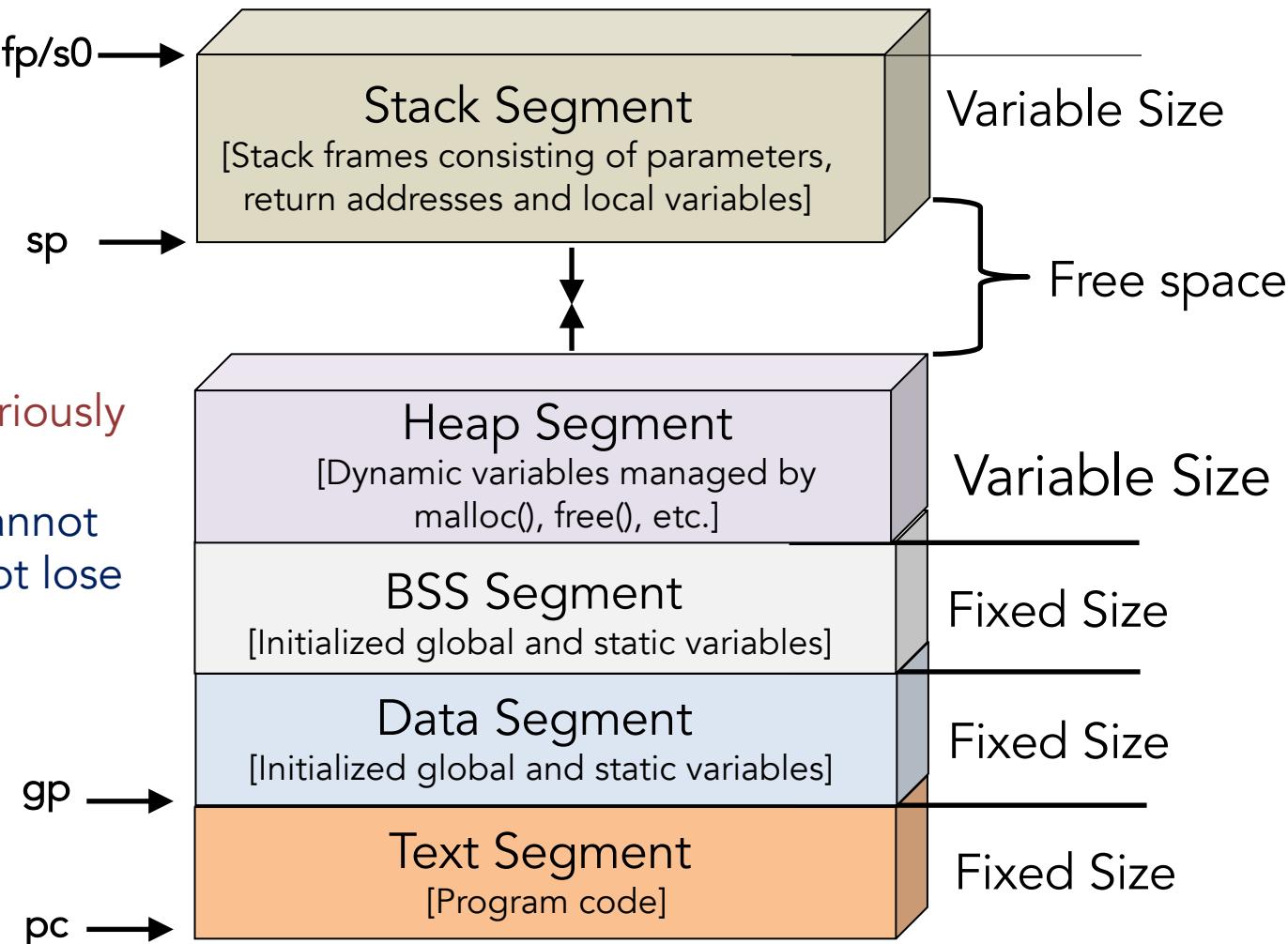
| Name | Reg Number | Usage | Preserved across call? |
|---------|------------|---|------------------------|
| zero | x0 | The constant value 0 | Yes |
| ra | x1 | Return address | Yes |
| sp | x2 | Stack pointer | Yes |
| gp | x3 | Global pointer | Yes |
| tp | x4 | Thread pointer | Yes |
| t0 – t2 | x5 – x7 | Temporary registers | No |
| s0/fp | x8 | Frame pointer | Yes |
| s1 | x9 | Saved register | Yes |
| a0 - a1 | x10 – x11 | Function arguments/Return values | No |
| a2 – a7 | x12 – x17 | Function arguments | No |

Program memory management



Program memory management

- And albert gloriously declares:
 - I simply cannot and will not lose myself!!!



RISC-V Assembly

```

int    val1   = 5;
int    val2;
int    val3   = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void)
{
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

val1:
    .word 5
val2:
    .zero 4
val3:
    .word 2
val4:
    .word 7
    .word 8
    .word 9
string1:
    .string "abc"
string2:
    .byte 97
    .byte 98
    .byte 99
array1:
    .zero 5
array2:
    .zero 8
array3:
    .zero 6

```

RISC-V Assembly

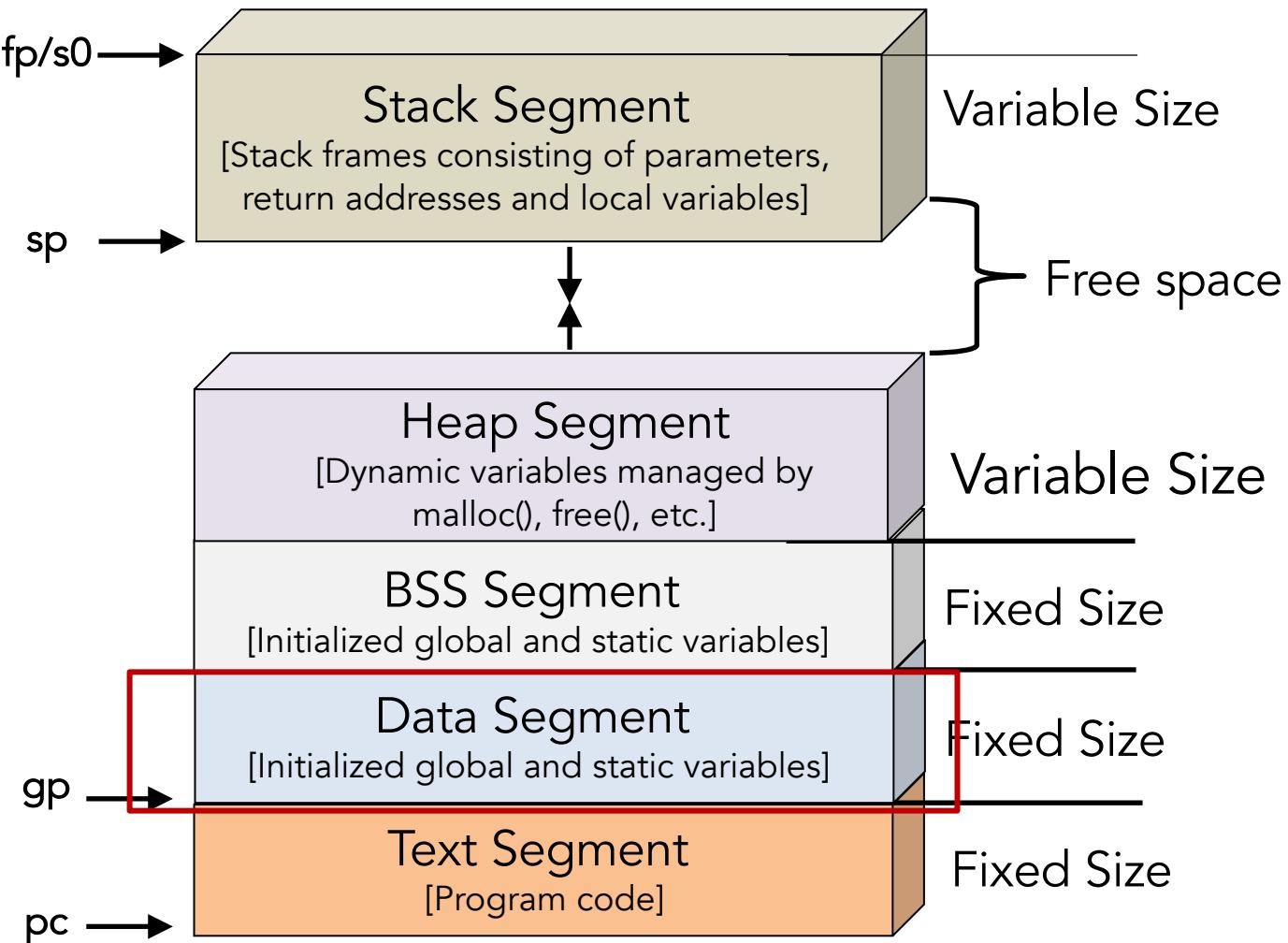
```

val1:
    .word 5
val2:
    .zero 4
val3:
    .word 2
val4:
    .word 7
    .word 8
    .word 9
string1:
    .string "abc"
string2:
    .byte 97
    .byte 98
    .byte 99
array1:
    .zero 5
array2:
    .zero 8
array3:
    .zero 6

```

| Address | Value | Address | Value |
|---------|-------|---------|-------|
| 0 | 5 | C | 7 |
| 1 | 0 | D | 0 |
| 2 | 0 | E | 0 |
| 3 | 0 | F | 0 |
| 4 | ? | 10 | 8 |
| 5 | ? | 11 | 0 |
| 6 | ? | 12 | 0 |
| 7 | ? | 13 | 0 |
| 8 | 2 | 14 | 9 |
| 9 | 0 | 15 | 0 |
| A | 0 | 16 | 0 |
| B | 0 | 17 | 0 |

Program memory management



RISC-V Assembly

```

val1:
    .word 5
val2:
    .zero 4
val3:
    .word 2
val4:
    .word 7
    .word 8
    .word 9
string1:
    .string "abc"
string2:
    .byte 97
    .byte 98
    .byte 99
array1:
    .zero 5
array2:
    .zero 8
array3:
    .zero 6

```

| Address | Value | Address | Value |
|---------|-------|---------|-------|
| 0 | 5 | C | 7 |
| 1 | 0 | D | 0 |
| 2 | 0 | E | 0 |
| 3 | 0 | F | 0 |
| 4 | ? | 10 | 8 |
| 5 | ? | 11 | 0 |
| 6 | ? | 12 | 0 |
| 7 | ? | 13 | 0 |
| 8 | 2 | 14 | 9 |
| 9 | 0 | 15 | 0 |
| A | 0 | 16 | 0 |
| B | 0 | 17 | 0 |

RISC-V Assembly

```

val1:
    .word 5
val2:
    .zero 4
val3:
    .word 2
val4:
    .word 7
    .word 8
    .word 9
string1:
    .string "abc"
string2:
    .byte 97
    .byte 98
    .byte 99
array1:
    .zero 5
array2:
    .zero 8

```

a Program code addresses vs. hardware addresses

| Address | Value | Address | Value |
|---------|-------|---------|-------|
| 8000 | 5 | 800C | 7 |
| 8001 | 0 | 800D | 0 |
| 8002 | 0 | 800E | 0 |
| 8003 | 0 | 800F | 0 |
| 8004 | ? | 8010 | 8 |
| 8005 | ? | 8011 | 0 |
| 8006 | ? | 8012 | 0 |
| 8007 | ? | 8013 | 0 |
| 8008 | 2 | 8014 | 9 |
| 8009 | 0 | 8015 | 0 |
| 800A | 0 | 8016 | 0 |
| 800B | 0 | 8017 | 0 |

RISC-V Assembly

```

val1:
    .word 5
val2:
    .zero 4
val3:
    .word 2
val4:
    .word 7
    .word 8
    .word 9
string1:
    .string "abc"
string2:
    .byte 97
    .byte 98
    .byte 99
array1:
    .zero 5
array2:
    .zero 8

```

Program code addresses vs. hardware addresses

| Address | Value |
|---------|-------|
| 8018 | 0x61 |
| 8019 | 0x62 |
| 801A | 0x63 |
| 801B | 0 |
| 801C | 0x61 |
| 801D | 0x62 |
| 801E | 0x63 |
| 801F | ? |
| 8020 | ? |
| 8021 | ? |
| 8022 | ? |
| 8023 | ? |

RISC-V Assembly – Let's Go All In

```
int    val1   = 5;
int    val2;
int    val3   = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}
```

RISC-V Assembly – Let's Go All In

```

int    val1  = 5;
int    val2;
int    val3  = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

Note that: num1, num2, num3 and num4 are a0, a1, a2, and a3, respectively.

```

foo:
add t0,a0,zero      # move num1 into t0
add t1,a1,zero      # move num2 into t1
add t2,a2,zero      # move num3 into t2

add t2,t2,a3        # t2 = num1 + num2

add t0,t0,t1        # t0 = num3 + num4

sub t1,t2,t0
# result = (num1 + num2) - (num3 + num4);

add a0,t1,zero
# result returned in a0
jr ra

```

RISC-V Assembly – Let's Go All In

```

int    val1  = 5;
int    val2;
int    val3  = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

Silly Albert! Who are you kidding!

Note that: num1, num2, num3 and num4 are a0, a1, a2, and a3, respectively.

```

foo:
add t0,a0,zero      # move num1 into t0
add t1,a1,zero      # move num2 into t1
add t2,a2,zero      # move num3 into t2

```

```

add t2,t2,a3        # t2 = num1 + num2

add t0,t0,t1        # t0 = num3 + num4

sub t1,t2,t0
# result = (num1 + num2) - (num3 + num4);

add a0,t1,zero
# result returned in a0
jr ra

```

RISC-V Assembly – Let's Go All In

```

int    val1  = 5;
int    val2;
int    val3  = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.globl  val1
.section .sdata,"aw",@progbits
.align  2
.type   val1, @object
.size   val1, 4
val1:
.word   5

.globl  val2
.section .sbss,"aw",@nobits
.align  2
.type   val2, @object
.size   val2, 4
val2:
.zero   4

.globl  val3
.section .sdata
.align  2
.type   val3, @object
.size   val3, 4
val3:
.word   2

```

RISC-V Assembly – Let's Go All In

```

int    val1  = 5;
int    val2;
int    val3  = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.globl  val1
.section .sdata,"aw",@progbits
.align  2
.type   val1, @object
.size   val1, 4

val1:
.word   5

.globl  val2
.section .sbss,"aw",@nobits
.align  2
.type   val2, @object
.size   val2, 4

val2:
.zero   4

.globl  val3
.section .sdata
.align  2
.type   val3, @object
.size   val3, 4

val3:
.word   2

```

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.globl  val4
.data
.align  2
.type   val4, @object
.size   val4, 12
val4:
.word   7
.word   8
.word   9

.globl  string1
.section .sdata
.align  2
.type   string1, @object
.size   string1, 4
string1:
.string "abc"

.globl  string2
.align  2
.type   string2, @object
.size   string2, 3
string2:
.byte   97
.byte   98
.byte   99

```

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.globl  val4
.data
.align  2
.type   val4, @object
.size   val4, 12

val4:
.word   7
.word   8
.word   9

.globl  string1
.section .sdata
.align  2
.type   string1, @object
.size   string1, 4

string1:
.string "abc"

.globl  string2
.align  2
.type   string2, @object
.size   string2, 3

string2:
.byte   97
.byte   98
.byte   99

```

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.globl array1
.section .sbss
.align 2
.type  array1, @object
.size  array1, 5
array1:
.zero   5

.globl array2
.section .sbss
.align 2
.type  array2, @object
.size  array2, 8
array2:
.zero   8

.globl array3
.section .sbss
.align 2
.type  array3, @object
.size  array3, 6
array3:
.zero   6

```

RISC-V Assembly – Let's Go All In

```

int    val1  = 5;
int    val2;
int    val3  = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.text
.align 1
.globl foo
.type  foo, @function
foo:
    addi   sp,sp,-48
    sw    s0,44(sp)
    addi   s0,sp,48
    sw    a0,-36(s0)
    sw    a1,-40(s0)
    sw    a2,-44(s0)
    sw    a3,-48(s0)
    lw    a4,-36(s0)
    lw    a5,-40(s0)
    add   a4,a4,a5
    lw    a3,-44(s0)
    lw    a5,-48(s0)
    add   a5,a3,a5
    sub   a5,a4,a5
    sw    a5,-20(s0)
    lw    a5,-20(s0)
    mv    a0,a5
    lw    s0,44(sp)
    addi   sp,sp,48
    jr    ra

```

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b'} ■ Stack pointer adjustment

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.text
.align 1
.globl foo
.type  foo, @function
foo:
    addi   sp,sp,-48 1
    sw    s0,44(sp)
    addi  s0,sp,48
    sw    a0,-36(s0)
    sw    a1,-40(s0)
    sw    a2,-44(s0)
    sw    a3,-48(s0)
    lw    a4,-36(s0)
    lw    a5,-40(s0)
    add   a4,a4,a5
    lw    a3,-44(s0)
    lw    a5,-48(s0)
    add   a5,a3,a5
    sub   a5,a4,a5
    sw    a5,-20(s0)
    lw    a5,-20(s0)
    mv    a0,a5
    lw    s0,44(sp)
    addi  sp,sp,48
    jr    ra

```

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b'}      ▪ Save old frame pointer
                                ▪ Caller's frame pointer

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.text
.align 1
.globl foo
.type   foo, @function

foo:
    addi   sp,sp,-48
    sw    s0,44(sp)      2
    addi   s0,sp,48
    sw    a0,-36(s0)
    sw    a1,-40(s0)
    sw    a2,-44(s0)
    sw    a3,-48(s0)
    lw    a4,-36(s0)
    lw    a5,-40(s0)
    add   a4,a4,a5
    lw    a3,-44(s0)
    lw    a5,-48(s0)
    add   a5,a3,a5
    sub   a5,a4,a5
    sw    a5,-20(s0)
    lw    a5,-20(s0)
    mv    a0,a5
    lw    s0,44(sp)
    addi  sp,sp,48
    jr    ra

```

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[10];
char string2[10];

char array1[5];
int  array2[2];
short int arra[10];

int foo (int num1,
         int num2,
         int num3,
         int num4)
{
    int result,
        result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

- The frame pointer (fp/s0) points to the start of the stack frame
 - It points to the base of the stack frame
 - Parameters passed in to the subroutine remain at constant places relative to the frame pointer
- It does not move for the duration of the subroutine call

foo:

```

.text
.align 1
.globl foo
.type  foo, @function

addi   sp,sp,-48
sw    s0,44(sp)
addi   s0,sp,48
sw    a0,-36(s0)
sw    a1,-40(s0)
sw    a2,-44(s0)
sw    a3,-48(s0)
lw    a4,-36(s0)
lw    a5,-40(s0)
add   a4,a4,a5
lw    a3,-44(s0)
lw    a5,-48(s0)
add   a5,a3,a5
sub   a5,a4,a5
sw    a5,-20(s0)
lw    a5,-20(s0)
mv    a0,a5
lw    s0,44(sp)
addi   sp,sp,48
jr    ra

```

3

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2,
         int num3, int num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

- Save the function arguments
 - In case, something happens in the middle of the function execution
 - Like divide by zero

```

.text
.align 1
.globl foo
.type  foo, @function

foo:
    addi   sp,sp,-48
    sw    s0,44(sp)
    addi   s0,sp,48
    sw    a0,-36(s0)
    sw    a1,-40(s0)
    sw    a2,-44(s0)
    sw    a3,-48(s0)
    lw    a4,-36(s0)
    lw    a5,-40(s0)
    add   a4,a4,a5
    lw    a3,-44(s0)
    lw    a5,-48(s0)
    add   a5,a3,a5
    sub   a5,a4,a5
    sw    a5,-20(s0)
    lw    a5,-20(s0)
    mv    a0,a5
    lw    s0,44(sp)
    addi   sp,sp,48
    jr    ra

```

4

RISC-V Assembly – Let's Go All In

```

int val1 = 5;
int val2;
int val3 = 2;
int val4[3] = {7,8,9};

char string1[] = "abc";
char string2[3] = {'a','b'} • Perform actual operations

char array1[5];
int array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.text
.align 1
.globl foo
.type foo, @function
foo:
    addi    sp,sp,-48
    sw     s0,44(sp)
    addi   s0,sp,48
    sw    a0,-36(s0)
    sw    a1,-40(s0)
    sw    a2,-44(s0)
    sw    a3,-48(s0)
    lw     a4,-36(s0)
    lw     a5,-40(s0)
    add   a4,a4,a5
    lw    a3,-44(s0)
    lw    a5,-48(s0)
    add   a5,a3,a5
    sub   a5,a4,a5
    sw    a5,-20(s0)
    lw    a5,-20(s0)
    mv    a0,a5
    lw    s0,44(sp)
    addi  sp,sp,48
    jr    ra

```

5

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

▪ Return result

```

.text
.align 1
.globl foo
.type  foo, @function
foo:
    addi   sp,sp,-48
    sw    s0,44(sp)
    addi   s0,sp,48
    sw    a0,-36(s0)
    sw    a1,-40(s0)
    sw    a2,-44(s0)
    sw    a3,-48(s0)
    lw    a4,-36(s0)
    lw    a5,-40(s0)
    add   a4,a4,a5
    lw    a3,-44(s0)
    lw    a5,-48(s0)
    add   a5,a3,a5
    sub   a5,a4,a5
    sw    a5,-20(s0)
    lw    a5,-20(s0)
    mv    a0,a5
    lw    s0,44(sp)
    addi  sp,sp,48
    jr    ra

```

6

RISC-V Assembly – Let's Go All In

```

int  val1  = 5;
int  val2;
int  val3  = 2;
int  val4[3] = {7,8,9};

char string1[]  = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int  array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1,
}

```

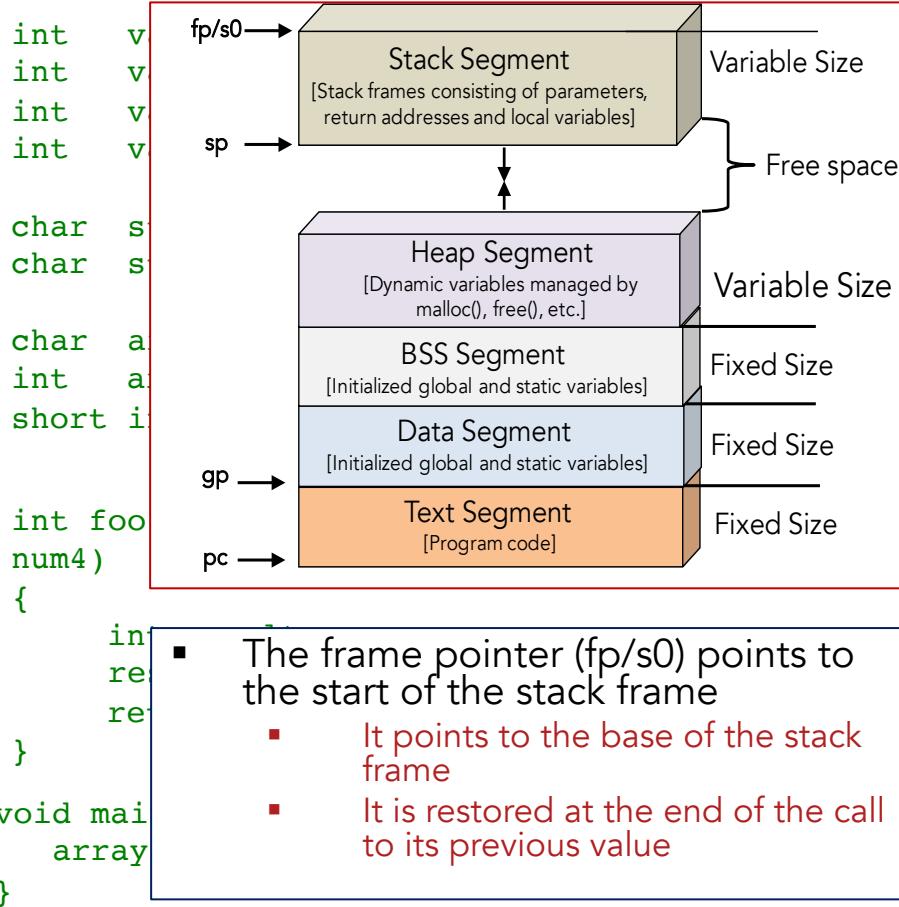
- Rollback things (frame pointer and stack pointer)

```

.text
.align 1
.globl foo
.type  foo, @function
foo:
    addi   sp,sp,-48
    sw    s0,44(sp)
    addi   s0,sp,48
    sw    a0,-36(s0)
    sw    a1,-40(s0)
    sw    a2,-44(s0)
    sw    a3,-48(s0)
    lw    a4,-36(s0)
    lw    a5,-40(s0)
    add   a4,a4,a5
    lw    a3,-44(s0)
    lw    a5,-48(s0)
    add   a5,a3,a5
    sub   a5,a4,a5
    sw    a5,-20(s0)
    lw    a5,-20(s0)
    mv    a0,a5
    lw    s0,44(sp)
    addi   sp,sp,48
    jr    ra

```

RISC-V Assembly – Let's Go All In



```

.text
.align 1
.globl foo
.type foo, @function

foo:
    addi sp,sp,-48
    sw s0,44(sp)
    addi s0,sp,48
    sw a0,-36(s0)
    sw a1,-40(s0)
    sw a2,-44(s0)
    sw a3,-48(s0)
    lw a4,-36(s0)
    lw a5,-40(s0)
    add a4,a4,a5
    lw a3,-44(s0)
    lw a5,-48(s0)
    add a5,a3,a5
    sub a5,a4,a5
    sw a5,-20(s0)
    lw a5,-20(s0)
    mv a0,a5
    lw s0,44(sp)
    addi sp,sp,48
    jr ra
  
```

7

RISC-V Assembly – Let's Go All In

```

int val1 = 5;
int val2;
int val3 = 2;
int val4[3] = {7,8,9};

char string1[] = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

- Jump out and time to go Albert!

```

.text
.align 1
.globl foo
.type foo, @function
foo:
    addi sp,sp,-48
    sw s0,44(sp)
    addi s0,sp,48
    sw a0,-36(s0)
    sw a1,-40(s0)
    sw a2,-44(s0)
    sw a3,-48(s0)
    lw a4,-36(s0)
    lw a5,-40(s0)
    add a4,a4,a5
    lw a3,-44(s0)
    lw a5,-48(s0)
    add a5,a3,a5
    sub a5,a4,a5
    sw a5,-20(s0)
    lw a5,-20(s0)
    mv a0,a5
    lw s0,44(sp)
    addi sp,sp,48
    jr ra

```

RISC-V Assembly – Let's Go All In

```

int val1 = 5;
int val2;
int val3 = 2;
int val4[3] = {7,8,9};

char string1[] = "abc";
char string2[3] = {'a','b','c'};

char array1[5];
int array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

- Jump out and time to go Albert!

```

.text
.align 1
.globl foo
.type foo, @function
foo:
    addi sp,sp,-48
    sw s0,44(sp)
    addi s0,sp,48
    sw a0,-36(s0)
    sw a1,-40(s0)
    sw a2,-44(s0)
    sw a3,-48(s0)
    lw a4,-36(s0)
    lw a5,-40(s0)
    add a4,a4,a5
    lw a3,-44(s0)
    lw a5,-48(s0)
    add a5,a3,a5
    sub a5,a4,a5
    sw a5,-20(s0)
    lw a5,-20(s0)
    mv a0,a5
    lw s0,44(sp)
    addi sp,sp,48
    jr ra

```

RISC-V Assembly – Let's Go All In

```

int    val1  = 5;
int    val2;
int    val3  = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.zero   6
.text
.align  1
.globl main
.type   main, @function
main:
    addi   sp,sp,-16
    sw     ra,12(sp)
    sw     s0,8(sp)
    addi   s0,sp,16
    lui    a5,%hi(val1)
    lw     a4,%lo(val1)(a5)
    lui    a5,%hi(val2)
    lw     a1,%lo(val2)(a5)
    lui    a5,%hi(val3)
    lw     a2,%lo(val3)(a5)
    lui    a5,%hi(val4)
    addi   a5,a5,%lo(val4)
    lw     a5,4(a5)
    mv    a3,a5
    mv    a0,a4
    call   foo
    mv    a4,a0
    lui    a5,%hi(array2)
    sw     a4,%lo(array2)(a5)
    li    a5,0
    mv    a0,a5
    lw     ra,12(sp)
    lw     s0,8(sp)
    addi   sp,sp,16
    jr    ra

```

RISC-V Assembly – Let's Go All In

```

int    val1  = 5;
int    val2;
int    val3  = 2;
int    val4[3] = {7,8,9};

char   string1[]  = "abc";
char   string2[3] = {'a','b','c'};

char   array1[5];
int    array2[2];
short int array3[3];

int foo (int num1, int num2, int num3, int
num4)
{
    int result;
    result = (num1 + num2) - (num3 + num4) ;
    return result;
}

void main(void) {
    array2 [0] = foo(val1, val2, val3, val4[1]);
}

```

```

.zero   6
.text
.align  1
.globl main
.type   main, @function

main:
    addi   sp,sp,-16
    sw     ra,12(sp)
    sw     s0,8(sp)
    addi   s0,sp,16
    lui    a5,%hi(val1)
    lw     a4,%lo(val1)(a5)
    lui    a5,%hi(val2)
    lw     a1,%lo(val2)(a5)
    lui    a5,%hi(val3)
    lw     a2,%lo(val3)(a5)
    lui    a5,%hi(val4)
    addi   a5,a5,%lo(val4)
    lw     a5,4(a5)
    mv    a3,a5
    mv    a0,a4
    call  foo
    mv    a4,a0
    lui    a5,%hi(array2)
    sw     a4,%lo(array2)(a5)
    li    a5,0
    mv    a0,a5
    lw     ra,12(sp)
    lw     s0,8(sp)
    addi   sp,sp,16
    jr    ra

```

Pseudo Instructions

| Useful Pseudo-instructions | Meaning | Actual Translation |
|--|--|--|
| <i>mv rd, rs</i> | Copy register rs to rd | addi rd, rs, 0 |
| <i>li rd, immediate</i> | Load immediate into rd | addi rd, \$0, immed |
| <i>beqz rs, offset</i> <i>bnez rs, offset</i> <i>bgtz rs, offset</i> | Branch if = zero Branch if != zero Branch if > zero | beq rs, x0, offset bne rs, x0, offset blt x0, rs, offset |
| <i>j offset</i> <i>jal offset</i> <i>jr rs</i> <i>jalr rs</i> | Jump Jump and link Jump register Jump and link register | jal x0, offset jal x1, offset jalr x0, rs, 0 jalr x1, rs, 0 |
| <i>call offset</i> | Call far-away subroutine | auipc x6, offset[31:12] jalr x1, x6, offset[11:0] |
| <i>ret</i> | Return from subroutine | jalr x0, x1, 0 |
| <i>nop</i> | Do nothing useful | add x0, x0, x0 |

Function Entry

```
prolog:  
    addi    sp, sp, -16    # -16 (frame size) depends on the amount of space needed  
    sw      ra, 12(sp)    # offset must be framesize-4  
    sw      fp, 8(sp)     # store the caller's fp  
    sw      gp, 4(sp)     # store caller's gp  
    mv      fp, sp        # fp is now the new function's frame base pointer  
  
    # ...  
    # The actual functional body of the function  
    # ...  
  
epilog:  
    lw      gp, 4(sp)  
    lw      fp, 8(sp)  
    addi   sp, sp, 16    # reverse the effect of first instruction  
    jr      ra
```

Assembly Programs

- If (condition) then do this

```
if ( i == j )
    i++ ;
j++;
```

Assuming t1 stores i and t2 stores j:

```
bne t1, t2, NEXT    # branch if !( i == j )
addi t1, t1, 1        # i++
NEXT: addi t2, t2, 1    # j++
```

Assembly Programs

- If (condition) then do this else do that

```
if ( i == j && i == k )  
    i++ ;  
else  
    j++;  
  
j = i + k ;
```

Assuming t0 stores i, t1 stores j, and t2 stores k:

```
bne t0, t1, ELSE # cond1: branch if !( i == j )  
bne t0, t2, ELSE # cond2: branch if !( i == k )  
addi t0, t0, 1    # if-body: i++  
j NEXT           # jump over else  
ELSE: addi t1, t1, 1 # else-body: j++  
NEXT: add t1, t0, t2 # j = i + k
```

Assembly Programs

- While (**condition**) then keep doing this

```
while ( i < j )
{
    k = k * 2;
    i++;
}
```

Assuming t0 stores i, t1 stores j, and t2 stores k:

```
Loop: bge t0, t1, DONE # branch if ! ( i < j )
      add t2, t2, t2    # k = k * 2
      addi t0, t0, 1     # i++
      j Loop             # jump back to top of loop
DONE:
```

Assembly Programs

- For (these many times) do this

```
for (i = 0; i < j; i++)  
{  
    k = k * 2;  
}
```

Assuming t0 stores i, t1 stores j, and t2 stores k:

```
addi t0, $0, 0      # i = 0  
Loop: bge t0, t1, DONE      # branch if ! ( i < j )  
      add t2, t2, t2      # k = k * 2  
      addi t0, t0, 1      # i++  
      j Loop            # jump back to top of loop  
DONE:
```

Assembly Programs

- For (these many times) do this

```
switch( i ) {  
    case 1:  
        i++;  
        break;  
    case 2:  
        i += 2;  
        break;  
    case 3:  
        i += 3;  
        break;  
}
```

Assuming t0 stores i and t1 stores temp:

```
addi t1, zero, 1      # case 1: set temp to 1  
bne t0, t1, CASE2    # false: branch to case 2  
j CASE1_BODY         # true: branch to case 1 body  
CASE2:  
    addi t1, zero, 2      # case 2: set temp to 2  
    bne t0, t1, CASE3    # false: branch to case 3  
    j CASE2_BODY         # true: branch to case 2 body  
CASE3:  
    addi t1, zero, 3      # case 3: set temp to 3  
    bne t0, t1, EXIT     # false: branch to exit  
    j CASE3_BODY         # true: branch to case 3 body  
CASE1_BODY:  
    addi t0, t0, 1        # case 1 body: i++  
    j EXIT                # break  
CASE2_BODY:  
    addi t0, t0, 2        # case 2 body: i += 2  
    j EXIT                # break  
CASE3_BODY:  
    addi t0, t0, 3        # case 3 body: i += 3  
    j EXIT                # break  
EXIT:
```

Assembly Programs

- Short program

```
clear(int array[], int size)
{
    int i;
    for (i = 0; i < size; i += 1)
        array[i] = 0;
}
```

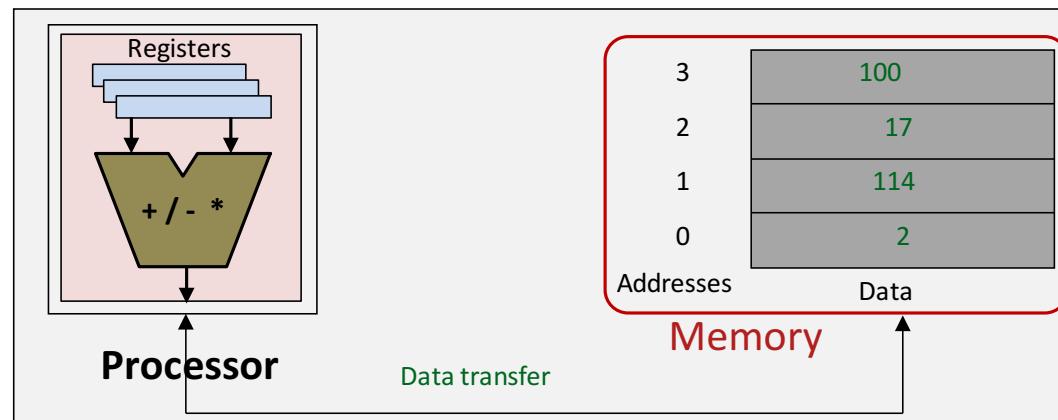
```
mv t0,zero      # i = 0

Loop: sll t1,t0,2      # t1 = i * 4
      add t2,a0,t1      # t2 = &array[i]
      sw zero, 0(t2)      # array[i] = 0
      addi t0,t0,1      # i = i + 1
      slt a3,t0,a1      # a3 = (i < size)
      bne a3,zero,Loop    # if (...)

      # goto loop
```

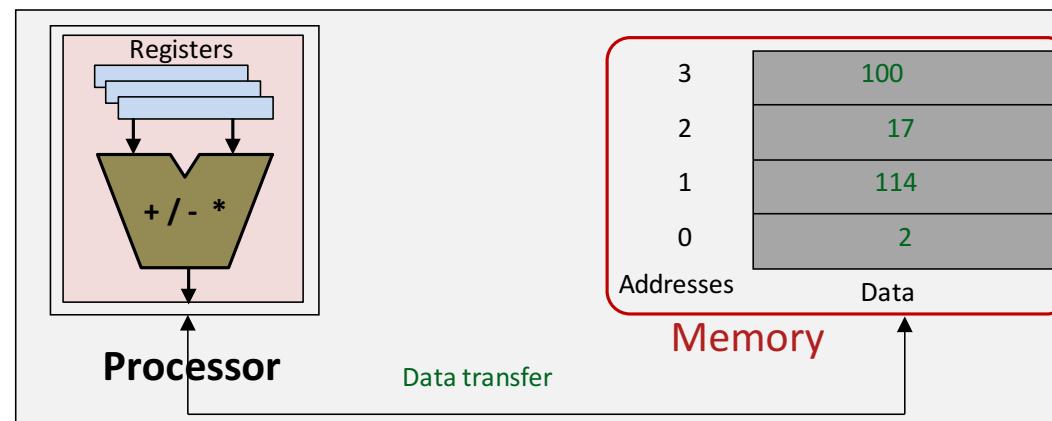
Data Addressing Modes

- RISC-V has four data addressing modes
 - 1. Register
 - Data is in a register
 - 2. Immediate
 - Data is specified in the instruction directly
 - 3. Displacement
 - Data is in memory with address calculated as Base + Offset
 - 4. PC-Relative
 - Data is in memory with address calculated as PC + Offset



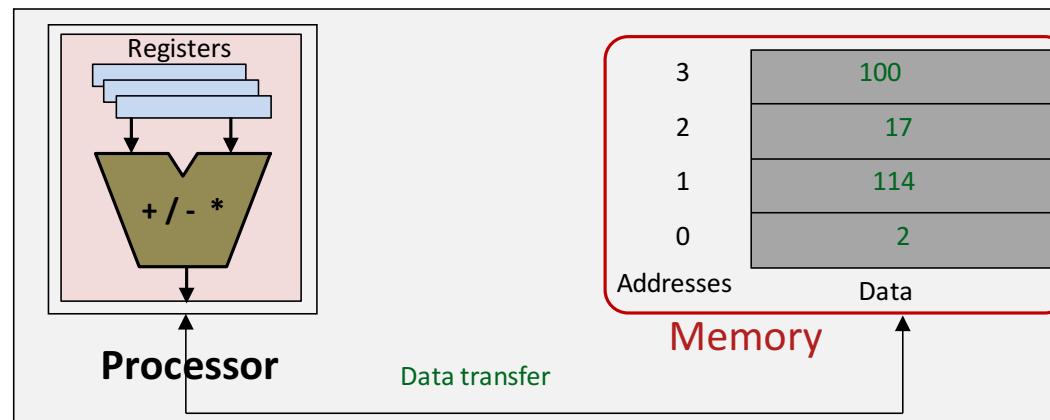
Data Addressing Modes

- RISC-V has four data addressing modes
 - Memory related addressing modes
 - Register indirect
 - $\text{add } x_2, (x_1)$
 - $x_1 \leftarrow x_2 + \text{Mem}[x_1]$
 - Displacement
 - $\text{add } x_2, 64(x_1)$
 - $x_1 \leftarrow x_2 + \text{Mem}[64+x_1]$

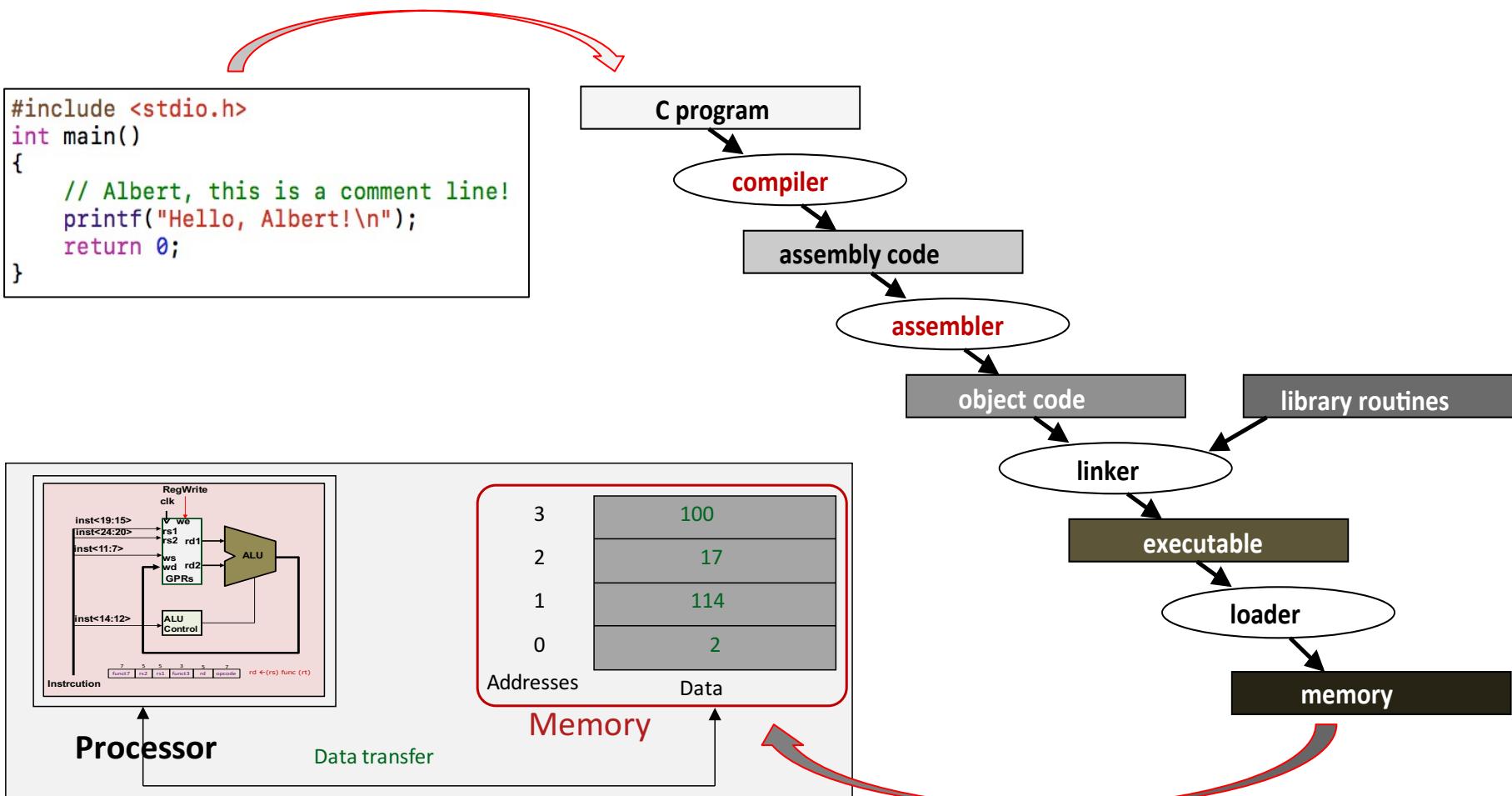


Data Addressing Modes

- Other memory address modes (non RISC-V)
 - Direct or absolute
 - $\text{add } x1, 1024$
 - $x1 \leftarrow x1 + \text{Mem}[1024]$
 - Memory indirect
 - $\text{add } x1, @x3$
 - $x1 \leftarrow x1 + \text{Mem}[\text{Mem}[x3]]$



The Big Picture



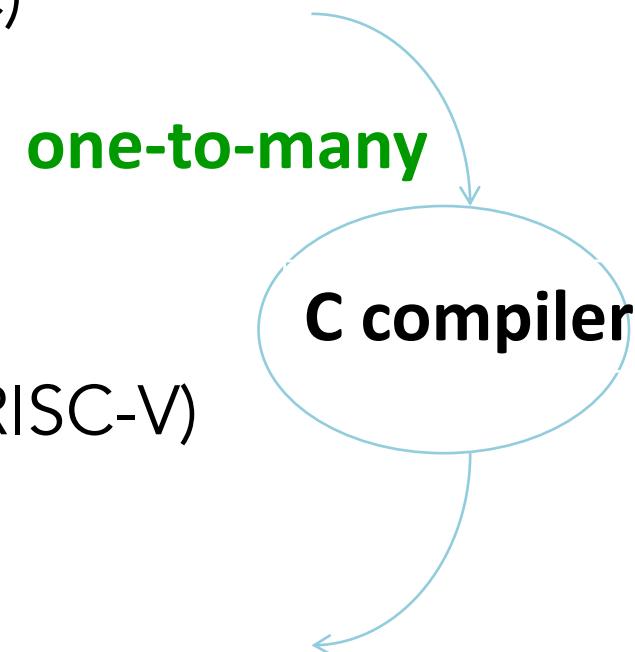
Application Compiling Process

- High-level language program (in C)

```
void swap ( int array[], int i) {  
    int temp;  
    temp      = array[i];  
    array[i]  = array[i+1];  
    array[i+1] = temp;
```

- Assembly language program (for RISC-V)
}

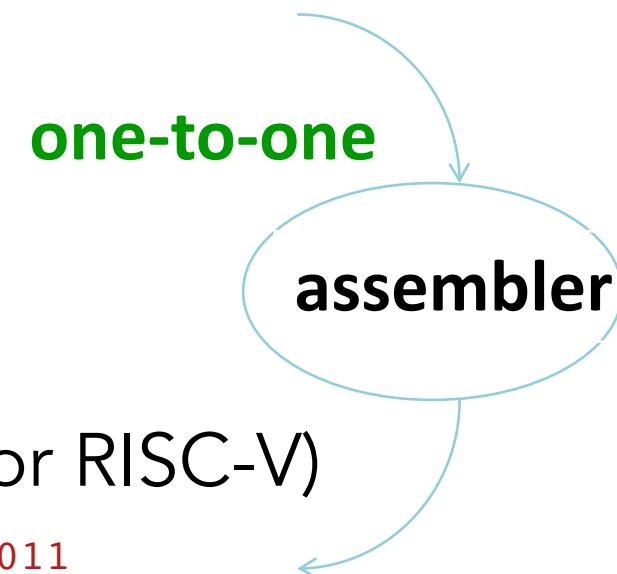
```
swap:  
    addi sp,sp,-48  
    ...  
    mv a5,a1  
    ...  
    ld s0,40(sp)  
    addi sp,sp,48  
    jr ra
```



Application Compiling Process

- Assembly language program (for RISC-V)

```
swap:  
    addi sp,sp,-48  
    ...  
    mv a5,a1  
    ...  
    ld s0,40(sp)  
    addi sp,sp,48  
    jr ra
```



- Machine (object, binary) code (for RISC-V)

| | | | |
|--------------|-------|-----------|---------|
| 111111010000 | 00010 | 000 00010 | 0010011 |
| 000000110000 | 00010 | 000 01000 | 0010011 |
| ... | | | |

Application Compiling Process

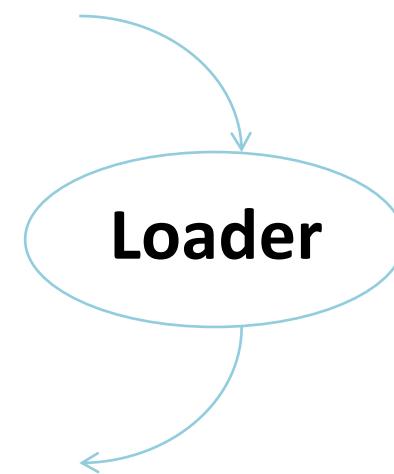
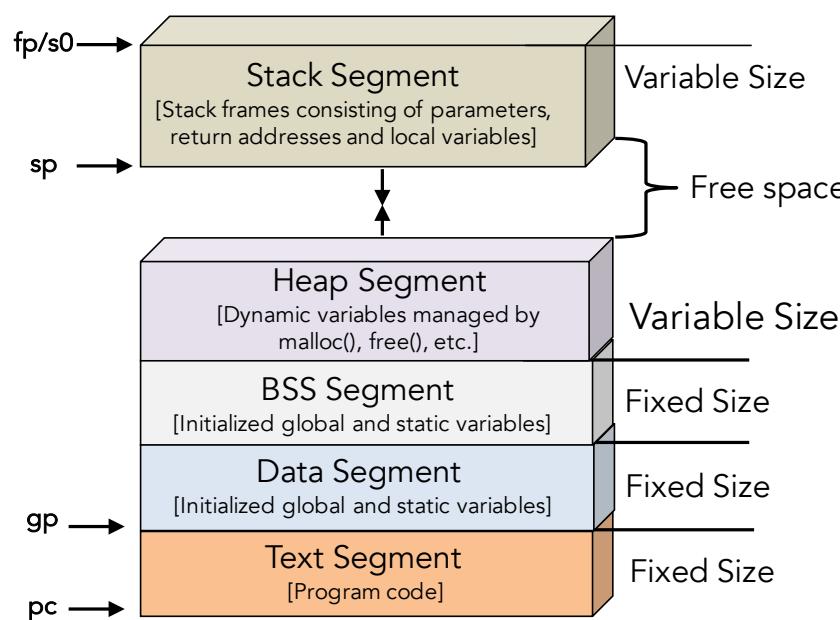
- Assembler (or compiler) translates program into machine instructions
- Provides information for building a complete program from the pieces
 - Header: described contents of object module
 - Text segment: translated instructions
 - Static data segment: data allocated for the life of the program
 - Relocation info: for contents that depend on absolute location of loaded program
 - Symbol table: global definitions and external refs
 - Debug info: for associating with source code

Application Compiling Process

- Machine (object, binary) code (for RISC-V)

```
111111010000 00010 000 00010 0010011
000000110000 00010 000 01000 0010011
```

- Memory organization



Application Compiling Process

- Load from image file on disk into memory
 1. Read header to determine segment sizes
 2. Create virtual address space
 3. Copy text and initialized data into memory
 4. Set up arguments on stack
 5. Initialize registers (including sp, fp, gp)
 6. Jump to startup routine
- Copies arguments to a0, ... and calls main

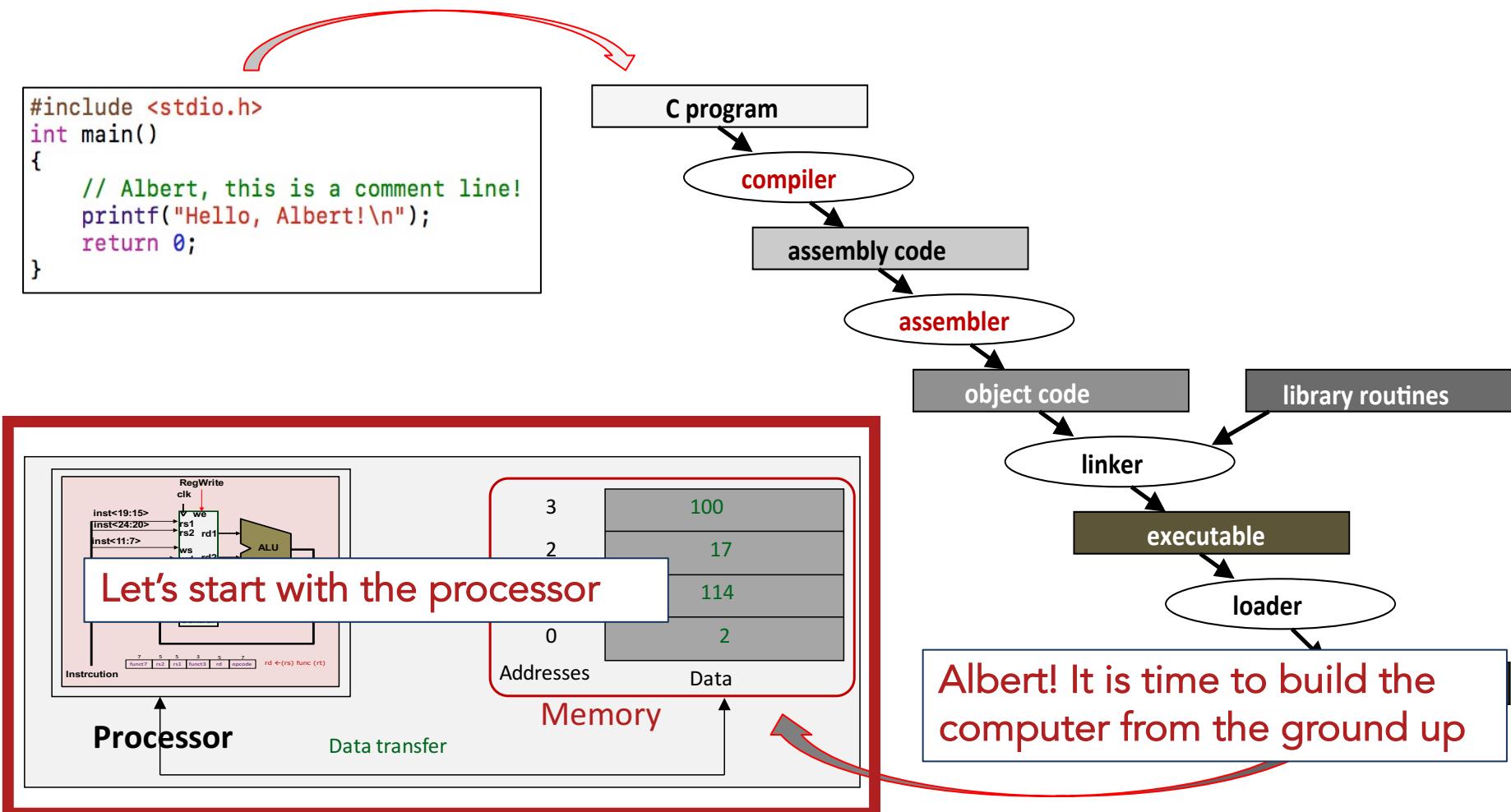
The Big Idea in Today's Computers

- Stored Program Computer

Program = A sequence of instructions

- Instructions represented in binary, just like data
- Instructions and data stored in memory
- Programs can operate on programs
 - e.g., compilers, linkers, ...
- Binary compatibility allows compiled programs to work on different computers
 - Standardized ISAs

The Big Picture



Next Lecture Module

- RISC-V ISA